- 1. Show that the set of all finite bitstrings  $\{0,1\}^*$  is countable.
- 2. Describe a Turing machine that computes function y = 2x + 1.
- 3. Show that  $3n^2 + 6n + 7 = O(n^2)$ .
- 4. Show that  $2n^3 + 6n^2 + 6n + 1 = O(n^3)$ .
- 5. Show that  $n^3 \neq O(n^2)$ .
- 6. Show that  $n! \neq O(2^n)$ .
- 7. Find functions f(n) and g(n) such that  $f(n) = O(g(n)), g(n) \neq O(f(n))$ , and  $f(n) \neq o(g(n))$ .
- 8. Given a list of functions in asymptotic notation, order them by growth rate (slowest to fastest).
- 9. Show that

10. Check if the following conditions are true

(a) 
$$\Theta(n+30) = \Theta(3n-1)$$
,  
(b)  $\Theta(n^2+2n-10) = \Theta(n^2+3n)$ ,  
(c)  $\Theta(n^3 \cdot 3n) = \Theta(n^2+3n)$ .

11. Write each of the following functions in O notation.

(a)  $5 + 0.001n^3 + 0.025n$  (b)  $500n + 100n^{1.5}$  (c)  $0.3n + 5n^{1.5} + 2.5n^{1.75}$ 

- 12. Reduce SAT to 3–coloring.
- 13. Reduce 3–coloring to SAT.
- 14. Reduce SAT to clique.
- 15. Reduce clique to independent set.
- 16. Reduce clique to vertex cover.
- 17. Reduce independent set to vertex cover.
- 18. Reduce DDHP to DLP.
- 19. Reduce CDHP to DLP.
- 20. Reduce DDHP to CDHP.

21. Is graph G shown in Fig. 2 3–colorable? In case it is 3–colorable, provide a valid 3-coloring as a proof. If it is not 3-colorable, how could we prove this fact?

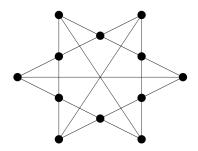


Figure 1: Graph G

22. Find a clique in the graph G below.

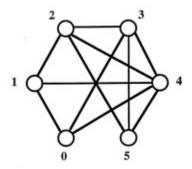


Figure 2: Graph G

## **1** Problem Definitions

**Definition 1** (Boolean satisfiability problem (SAT)). The problem of determining if there exists an interpretation that satisfies a given Boolean formula.

**Definition 2** (*k*-colorability problem). Given a graph, decide if it can be colored using k colors such that no two adjacent vertices are colored with the same color.

**Definition 3** (Clique problem). The problem of determining the existence of a complete subgraph (a subset of vertices all adjacent to each other) in a given graph.

**Definition 4** (Independent set problem). The problem of determining the existence of a subset of vertices of a graph, such that no two vertices are adjacent in this subset.

**Definition 5** (Vertex cover problem). The problem of determining the existence of a subset of vertices of a graph that includes at least one endpoint of every edge of the graph.

**Definition 6** (Discrete logarithm problem (DLP)). Given a multiplicative cyclic group G, its generator g, an element  $h \in G$ , find k such that  $g^k = h$  in G.

**Definition 7** (Computational Diffie-Hellman problm (CDHP)). Given a triplet  $(g, g^a, g^b)$  of elements of a multiplicative cyclic group G generated by g, for  $a, b \in Z_n$ , n being the number of elements in G, compute  $g^{ab}$  in G.

**Definition 8** (Decisional Diffie–Hellman problem (DDHP)). Given a triplet  $(g^a, g^b, g^c)$  of elements of a multiplicative cyclic group G generated by g, decide whether  $g^c = g^{ab}$  in G.